# The union-find problem 

Bruce Merry

## Background



## Simple solution 1



| 0 | 1 | 2 | 0 | 0 |
| :--- | :--- | :--- | :--- | :--- |

## Simple solution 2


$0: \quad \mathrm{A} \rightarrow \mathrm{D} \rightarrow \mathrm{E}$
1: B
2: C


## Trees



## Trees



## Pseudo-code

$$
\text { proc find }(x) \equiv
$$

$$
\underline{\text { while }} \operatorname{parent}(x) \geq 0
$$

$$
x:=\operatorname{parent}(x)
$$

## end

proc union $(x, y) \equiv$

$$
\begin{aligned}
& r_{x}:=\text { find }(x) \\
& r_{y}:=\text { find }(y) \\
& \text { if } r_{x} \neq r_{y} \\
& \quad \text { parent }\left(r_{y}\right):=r_{x}
\end{aligned}
$$

$$
\underline{\mathrm{fi}}
$$

## Balancing

Prevent degenerate trees:

- Make shallower tree the child of the deeper
- Make smaller tree the child of the larger Height/size can be stored in root.


## Path compression



## Path compression



## Minimum spanning trees



## Kruskal's algorithm

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- Repeat until $V-1$ edges have been added.

